

Secure Software with Formal Guarantees

Using Hax

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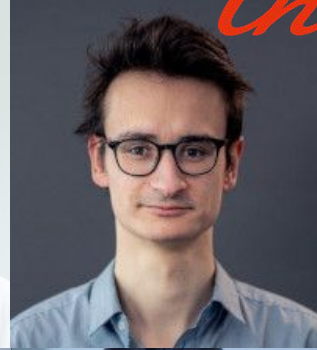
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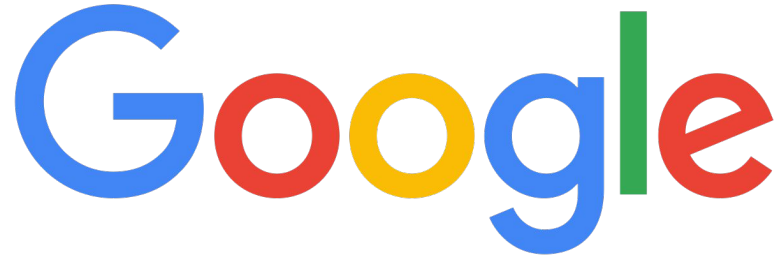
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Building Secure Software

[...] testing is a necessary but insufficient step in the development process to fully reduce vulnerabilities at scale [...]



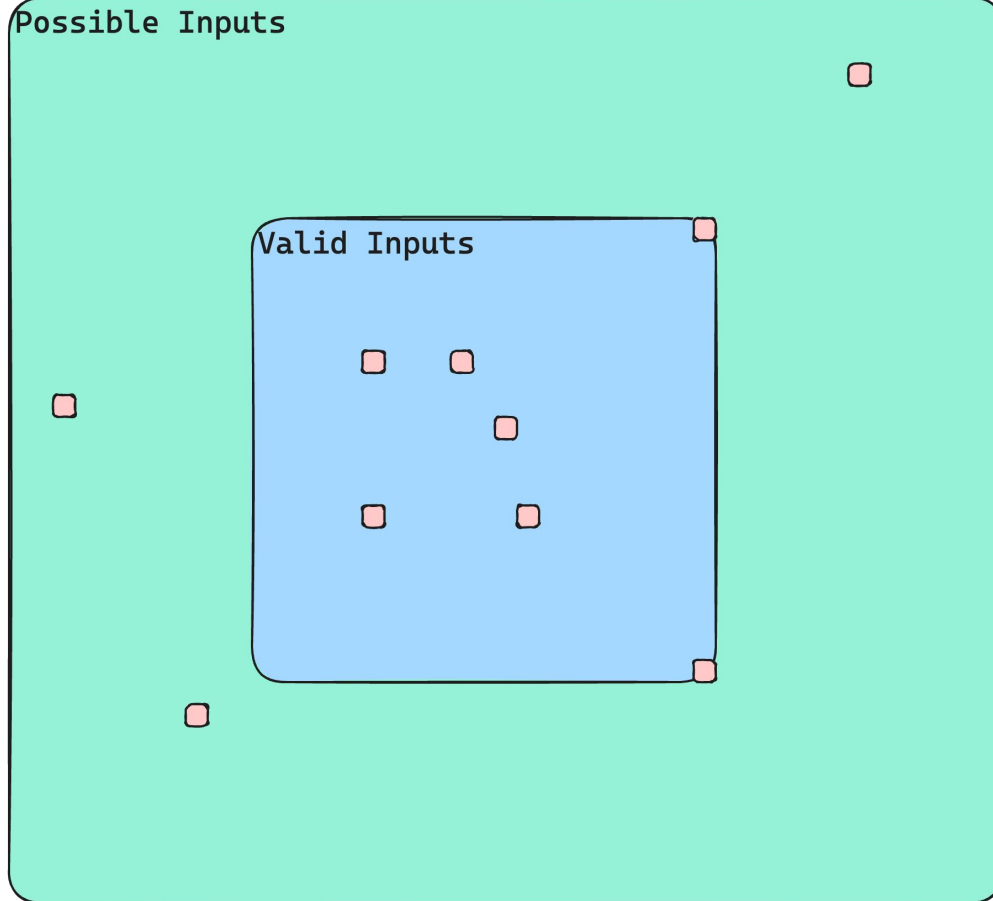
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Possible Inputs



Valid Inputs

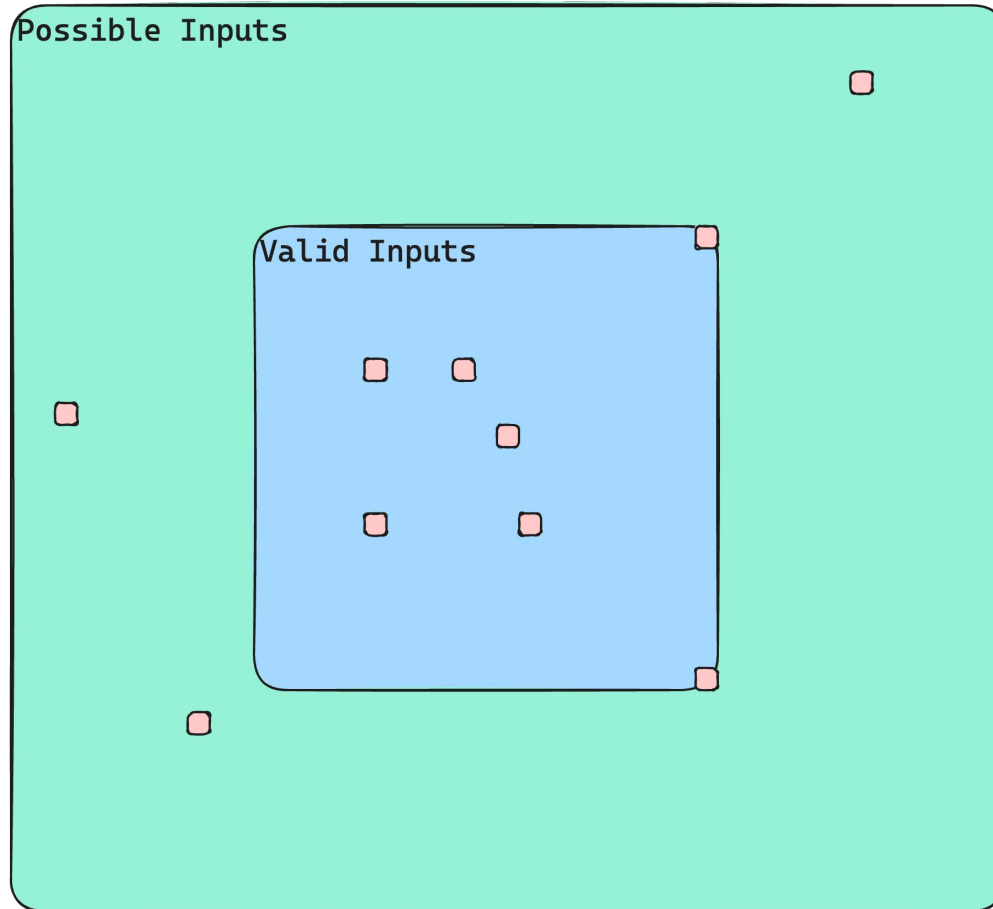
Testing



Testing

Wycheproof
ECDSA P256

471 Tests



Possible
Inputs

64 bytes
Signature

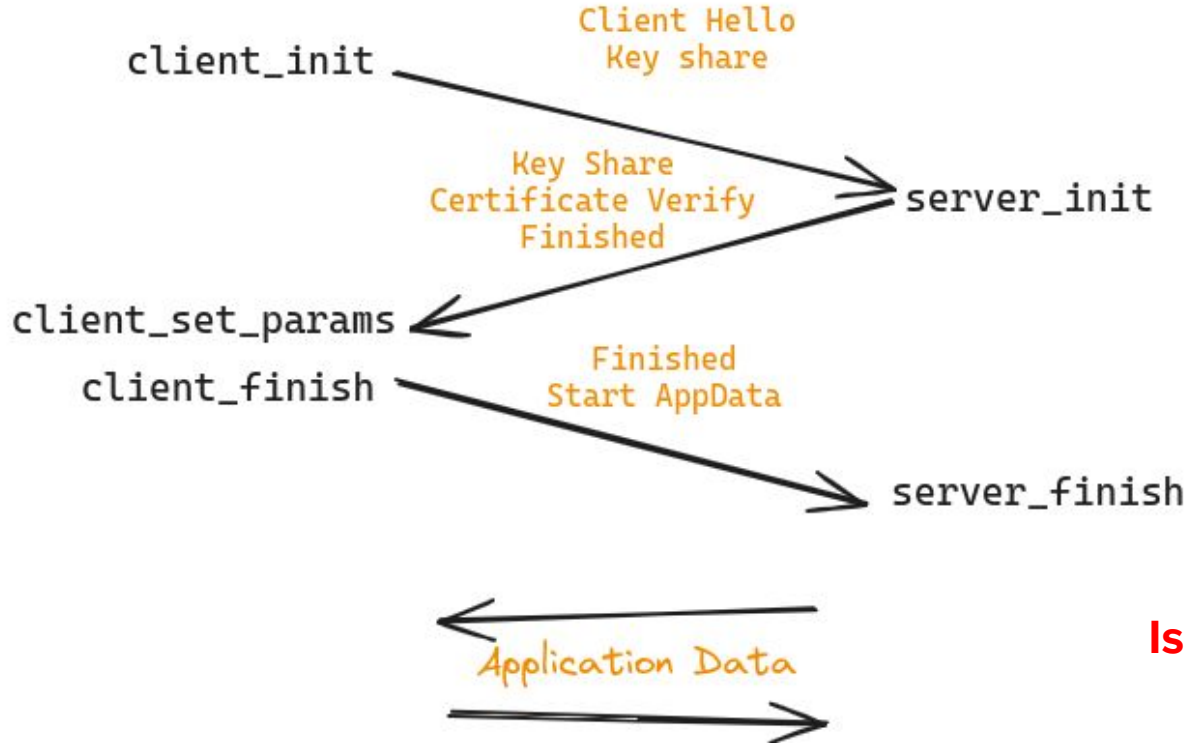
64 bytes
public key

Verification

Possible Inputs

Valid Inputs

TLS



Is this secure?

**“[...] correctness is defined as the ability
of a piece of software to meet a specific
[...] requirement”**



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Usable Verification Tools

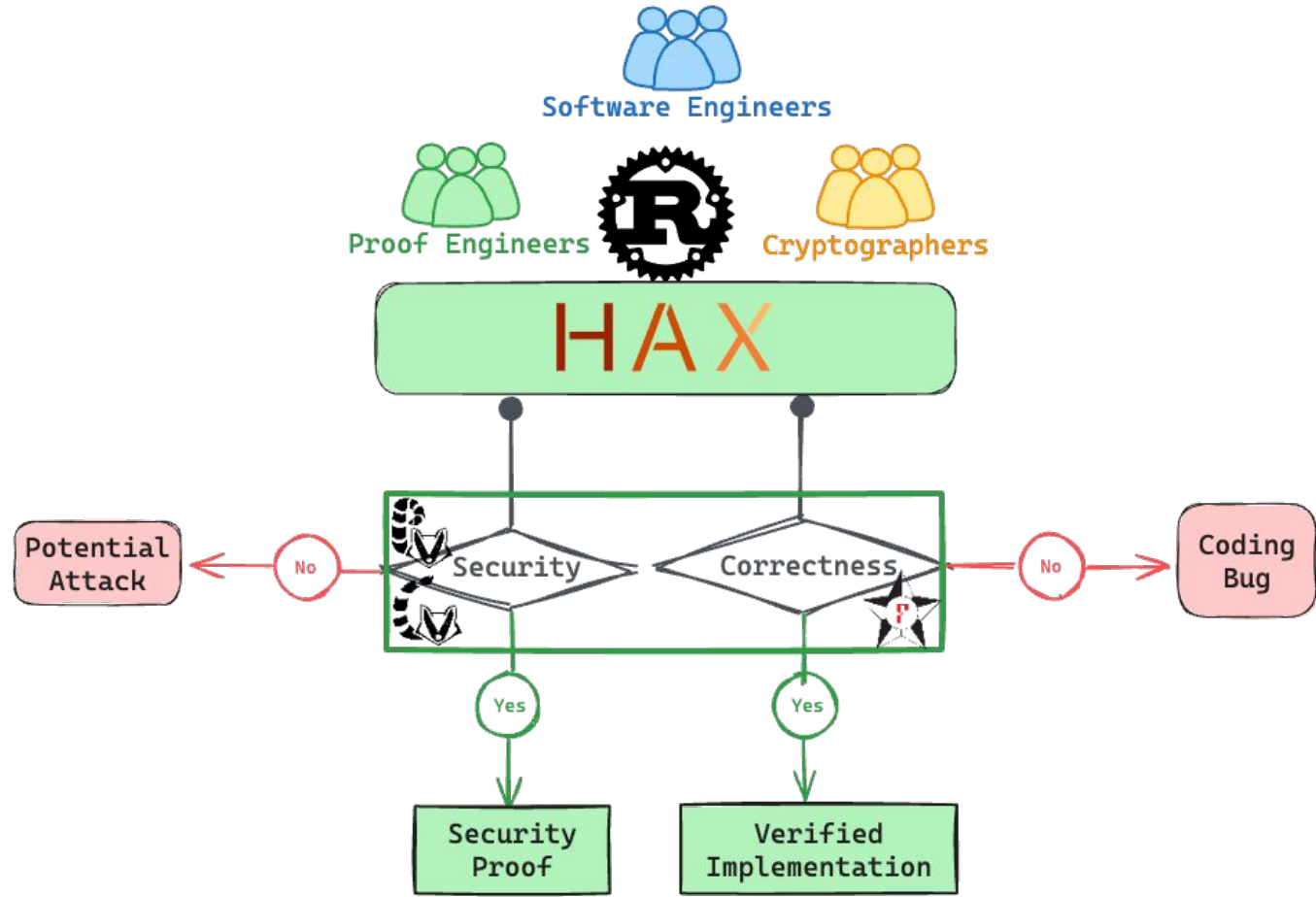
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Formal Verification

Correctness

HAX

verification
toolchain

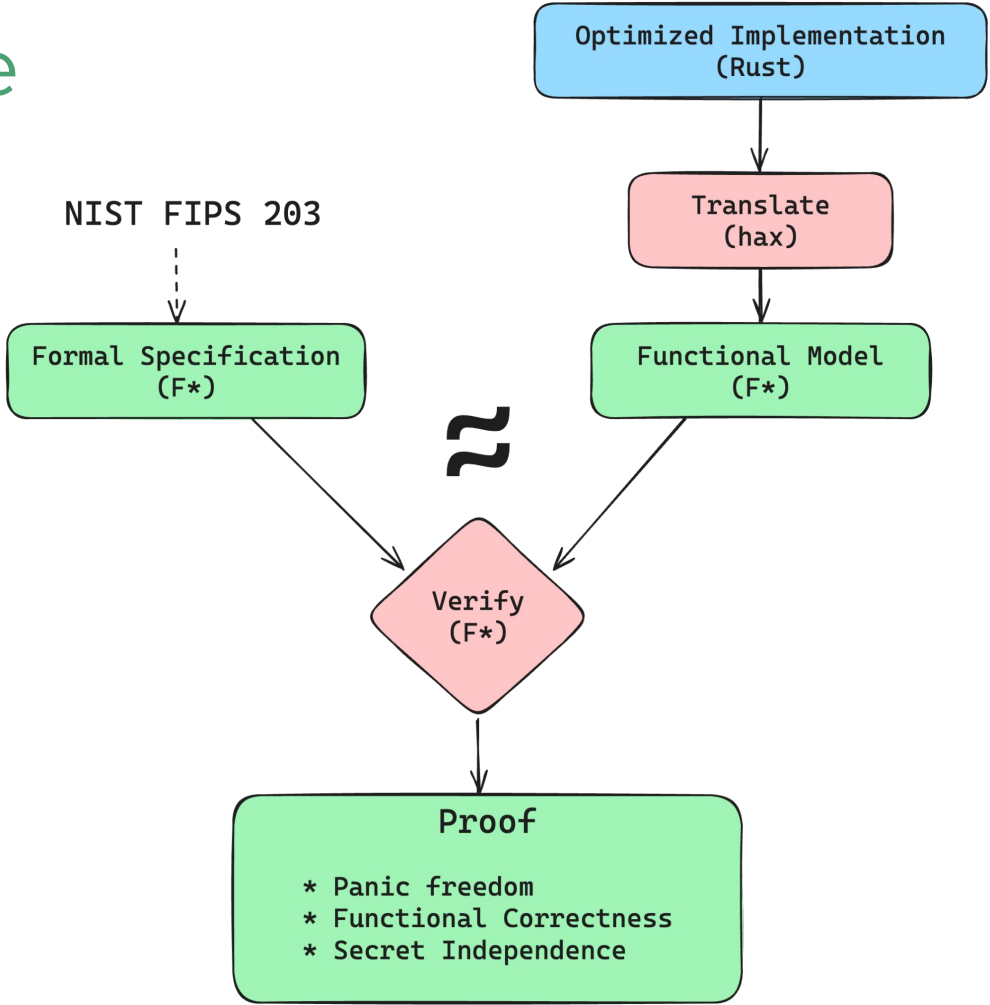


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Verifying Rust Code with hax and F*



The `hax` process

hax: Process

```
// reduce_once reduces  $0 \leq x < 2 * kPrime$ , mod kPrime.
static uint16_t reduce_once(uint16_t x) {
    assert(x < 2 * kPrime);
    const uint16_t subtracted = x - kPrime;
    uint16_t mask = 0u - (subtracted >> 15);
    // On Aarch64, omitting a |value_barrier_u16| results in a 2x speedup of Kyber
    // overall and Clang still produces constant-time code using `csel`. On other
    // platforms & compilers on godbolt that we care about, this code also
    // produces constant-time output.
    return (mask & x) | (~mask & subtracted);
}
```

```
// constant time reduce x mod kPrime using Barrett reduction. x must be less
// than kPrime + 2*kPrime2.
```

```
static uint16_t reduce(uint32_t x) {
    assert(x < kPrime + 2u * kPrime * kPrime);
    uint64_t product = (uint64_t)x * kBarrettMultiplier;
    uint32_t quotient = (uint32_t)(product >> kBarrettShift);
    uint32_t remainder = x - quotient * kPrime;
    return reduce_once(remainder);
}
```

hax: Process

```
// reduce_once reduces  $0 \leq x < 2 * kPrime$ , mod kPrime.
```

```
static uint16_t reduce_once(uint16_t x) {
```

```
    assert(x < 2 * kPrime);
```

```
    const uint16_t subtracted = x - kPrime;
```

```
    uint16_t mask = 0u - (subtracted >> 15);
```

```
    // On Aarch64, omitting a |value_barrier_u16| results in a 2x speedup of Kyber
```

```
    // overall and Clang still produces constant-time code using `csel`. On other
```

```
    // platforms & compilers on godbolt that we care about, this code also
```

```
    // produces constant-time output.
```

```
    return (mask & x) | (~mask & subtracted);
```

```
}
```

```
// constant time reduce x mod kPrime using Barrett reduction. x must be less
```

```
// than kPrime + 2*kPrime2.
```

```
static uint16_t reduce(uint32_t x) {
```

```
    assert(x < kPrime + 2u * kPrime * kPrime);
```

```
    uint64_t product = (uint64_t)x * kBarrettMultiplier;
```

```
    uint32_t quotient = (uint32_t)(product >> kBarrettShift);
```

```
    uint32_t remainder = x - quotient * kPrime;
```

```
    return reduce_once(remainder);
```

```
}
```

hax: Process

```
#[requires(coefficient_bits ≤ 11 && i32::from(fe) ≤ FIELD_MODULUS)]
#[ensures(|result| result ≥ 0 && result ≤ (1 << coefficient_bits) - 1)]
pub(super) fn compress_q(coefficient_bits: usize, fe: u16) → KyberFieldElement {
    let mut compressed: u32 = (fe as u32) << (coefficient_bits + 1);
    compressed += FIELD_MODULUS as u32;
    compressed ≐ (FIELD_MODULUS << 1) as u32;


    (compressed & ((1u32 << coefficient_bits) - 1)) as KyberFieldElement
}
```

1. Make the requirements formal
2. hax attributes for “design by contract”
3. F* statically checks that the properties hold

Example

Proving correctness
of Barrett reduction


Writing Crypto Code in Rust



```
pub(crate) fn barrett_reduce(input: i32) -> i32 {
    let t = (i64::from(input) * 20159) + (0x4_000_000 >> 1);
    let quotient = (t >> 26) as i32;
    let remainder = input - (quotient * 3329);
    remainder
}
```

Barrett Reduction: computes **input % 3329**
(in constant time)

Potential Panics in Rust Code



```
pub(crate) fn barrett_reduce(input: i32) -> i32 {  
    let t = (i64::from(input) * 20159) + (0x4_000_000 >> 1);  
    let quotient = (t >> 26) as i32;  
    let remainder = input - (quotient * 3329);  
    remainder  
}
```

These arithmetic operations may overflow or underflow causing the code to panic at run-time

Proving Panic Freedom and Correctness in F*

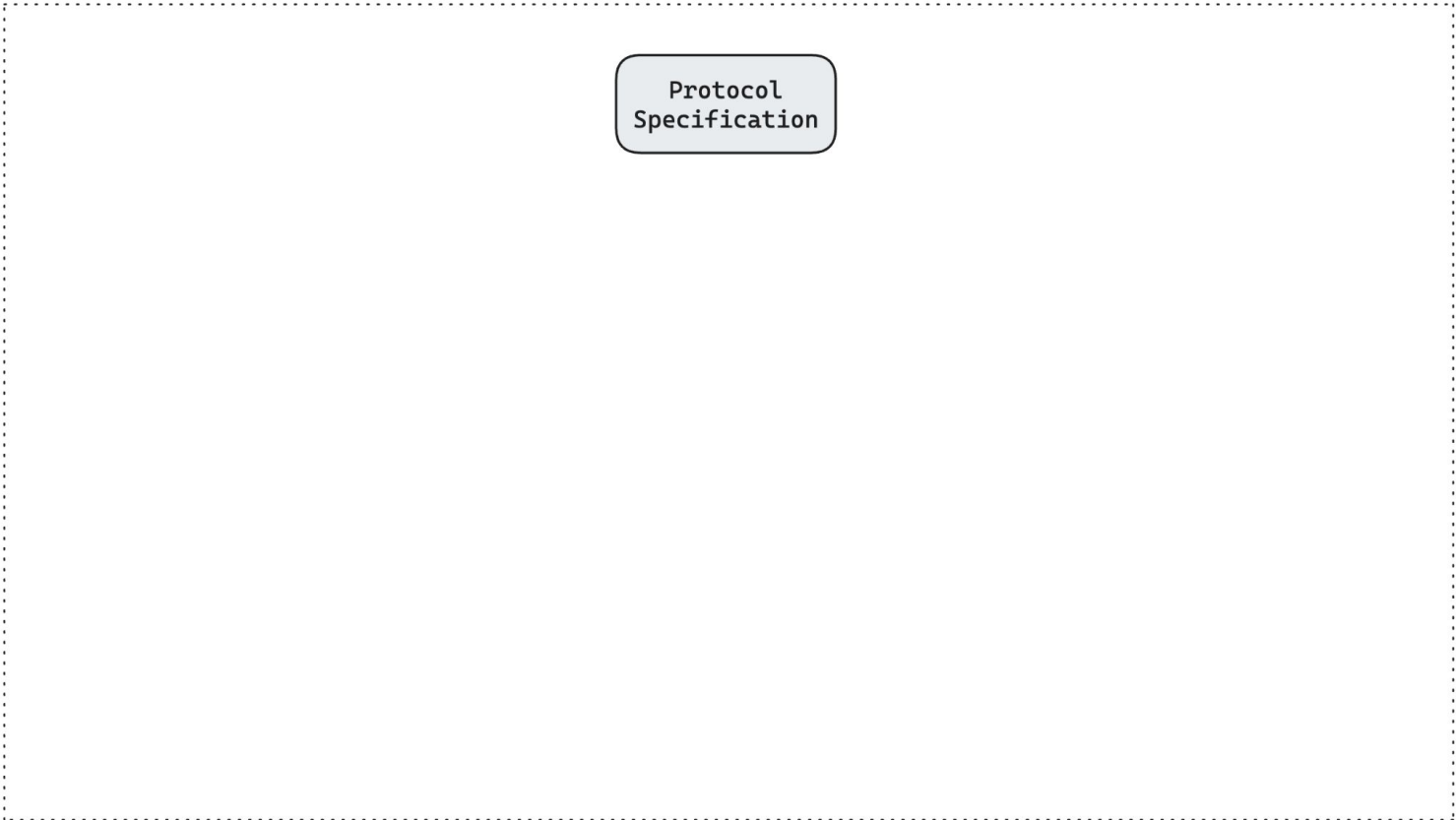
```
val barrett_reduce (input: i32_b (v v_BARRETT_R))
  : Pure (i32_b 3328)
  (requires True)
  (ensures fun result ->
    v result % v Libcrux.Kem.Kyber.Constants.v_FIELD_MODULUS
    = v input %v Libcrux.Kem.Kyber.Constants.v_FIELD_MODULUS)
```

Expected behaviour: $\text{result} \approx \text{input} \% 3329$

Formal Verification

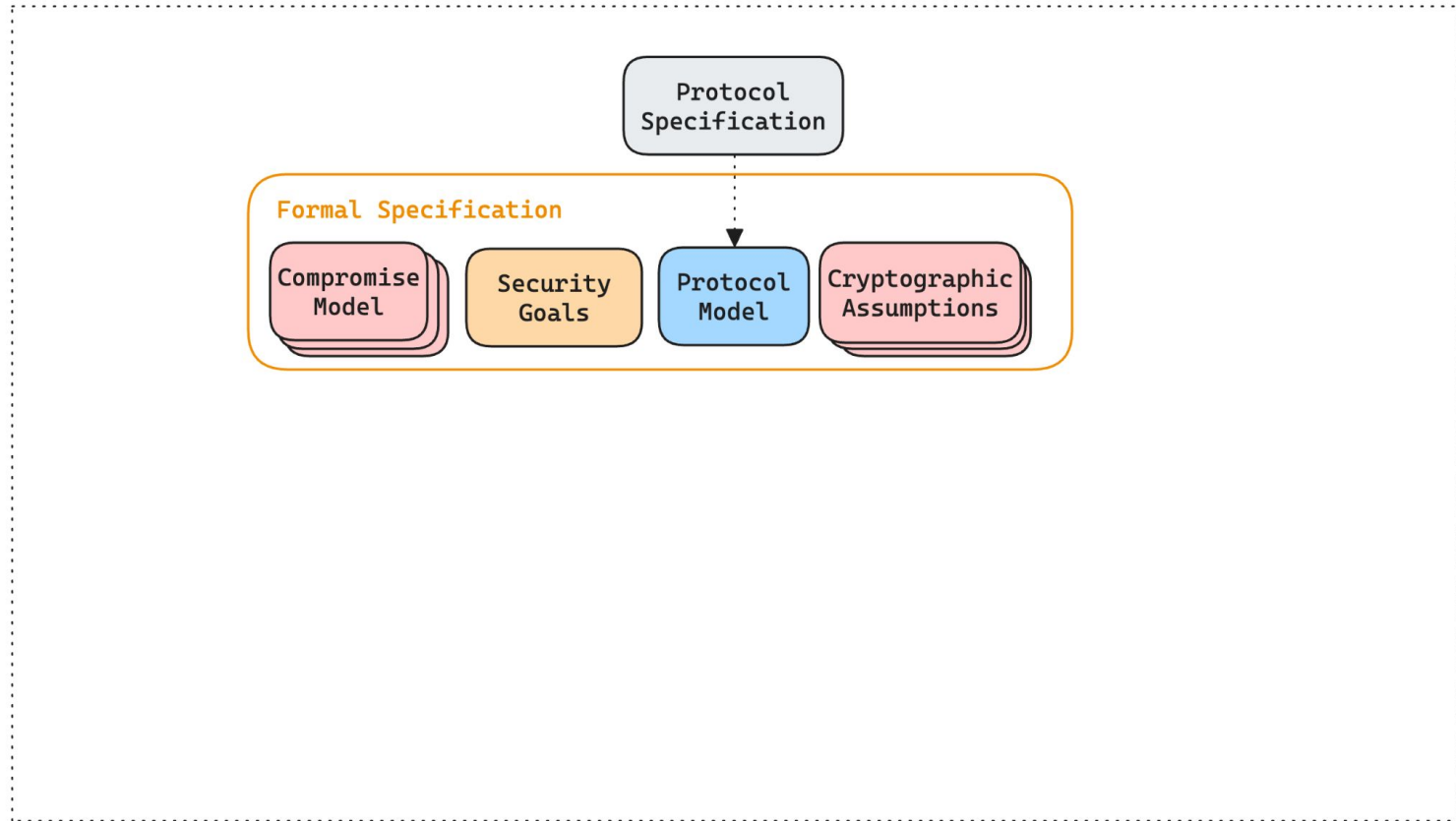
Security

Our Formal Verification Methodology | Security

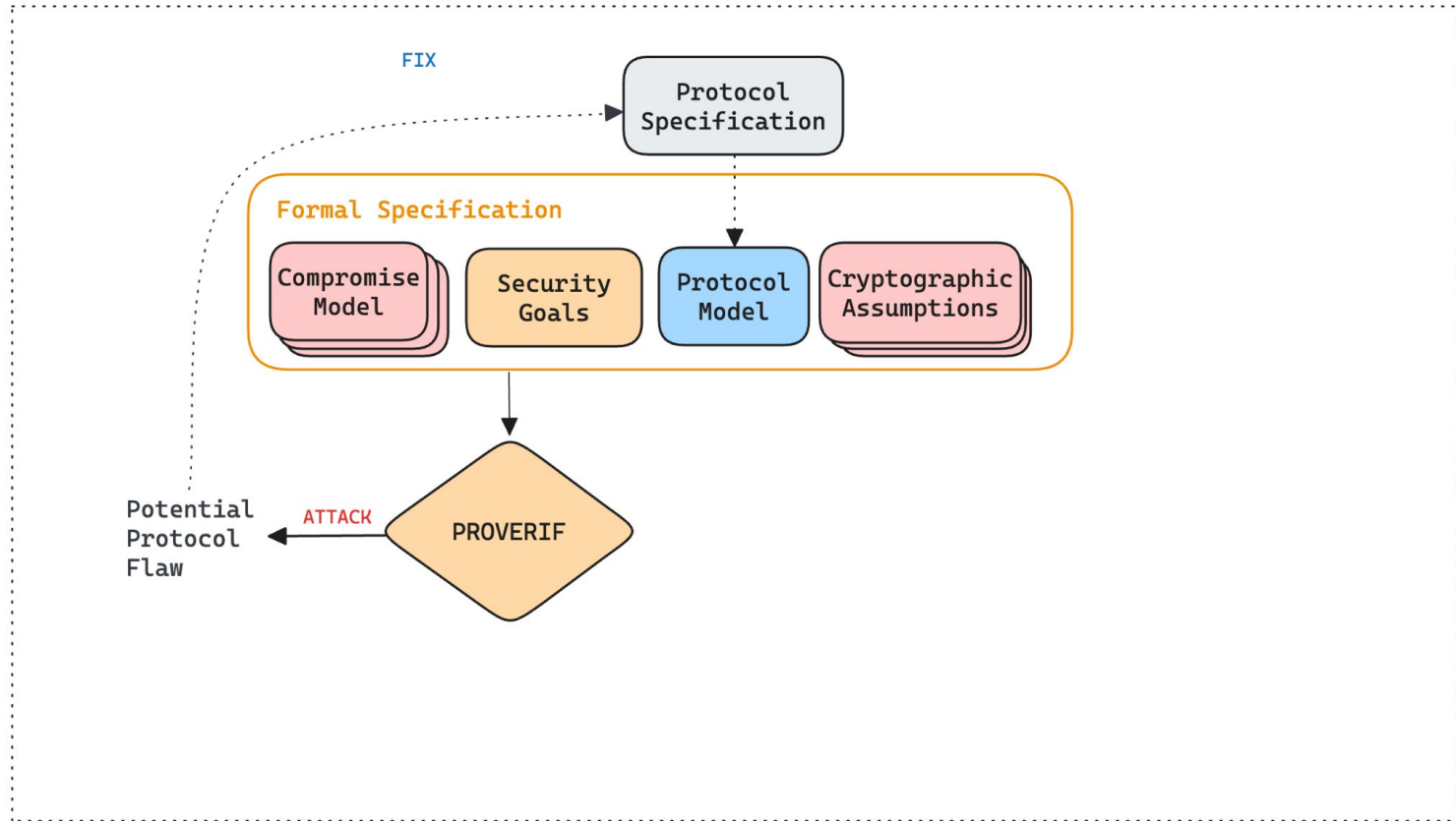


Protocol
Specification

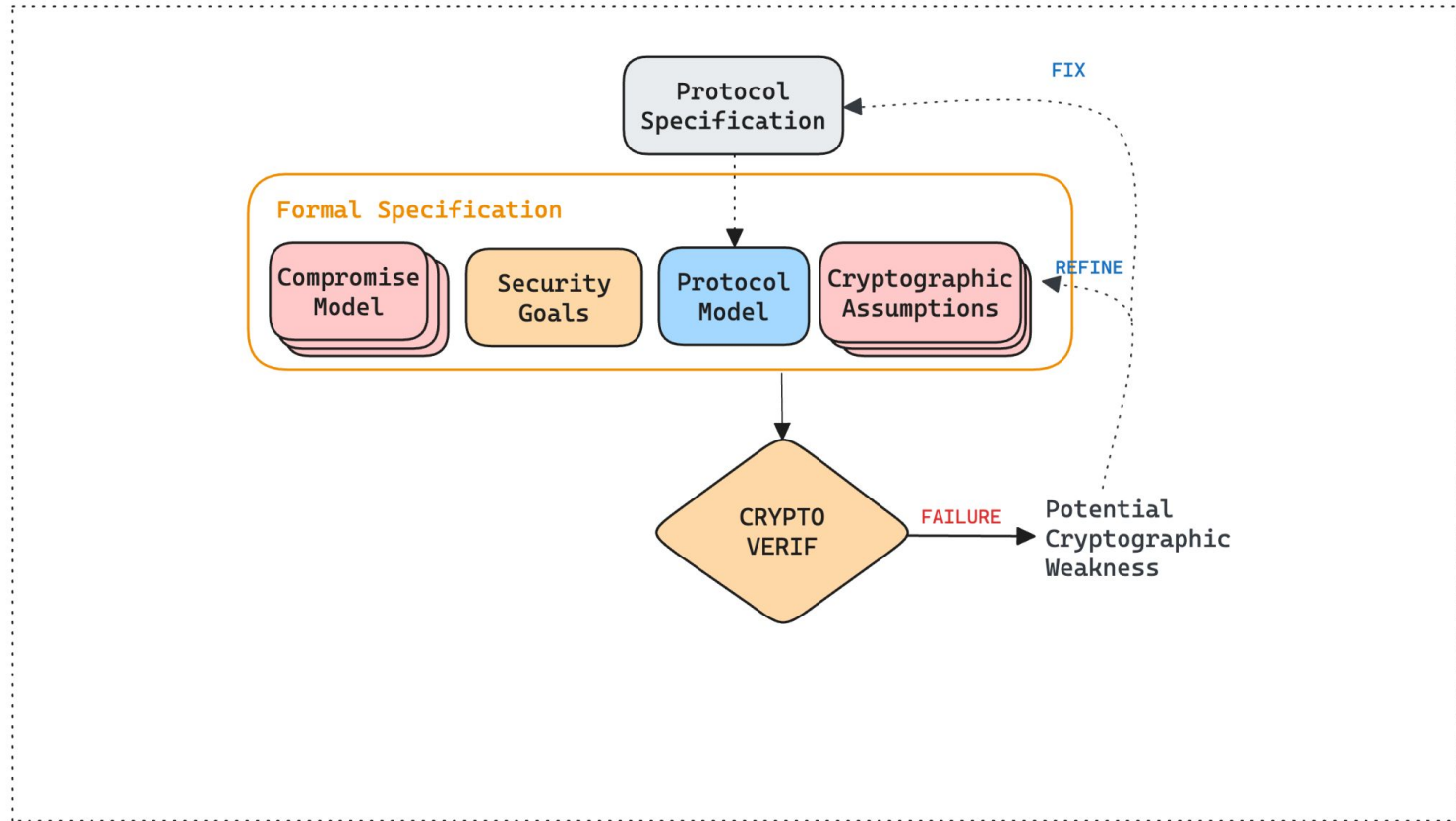
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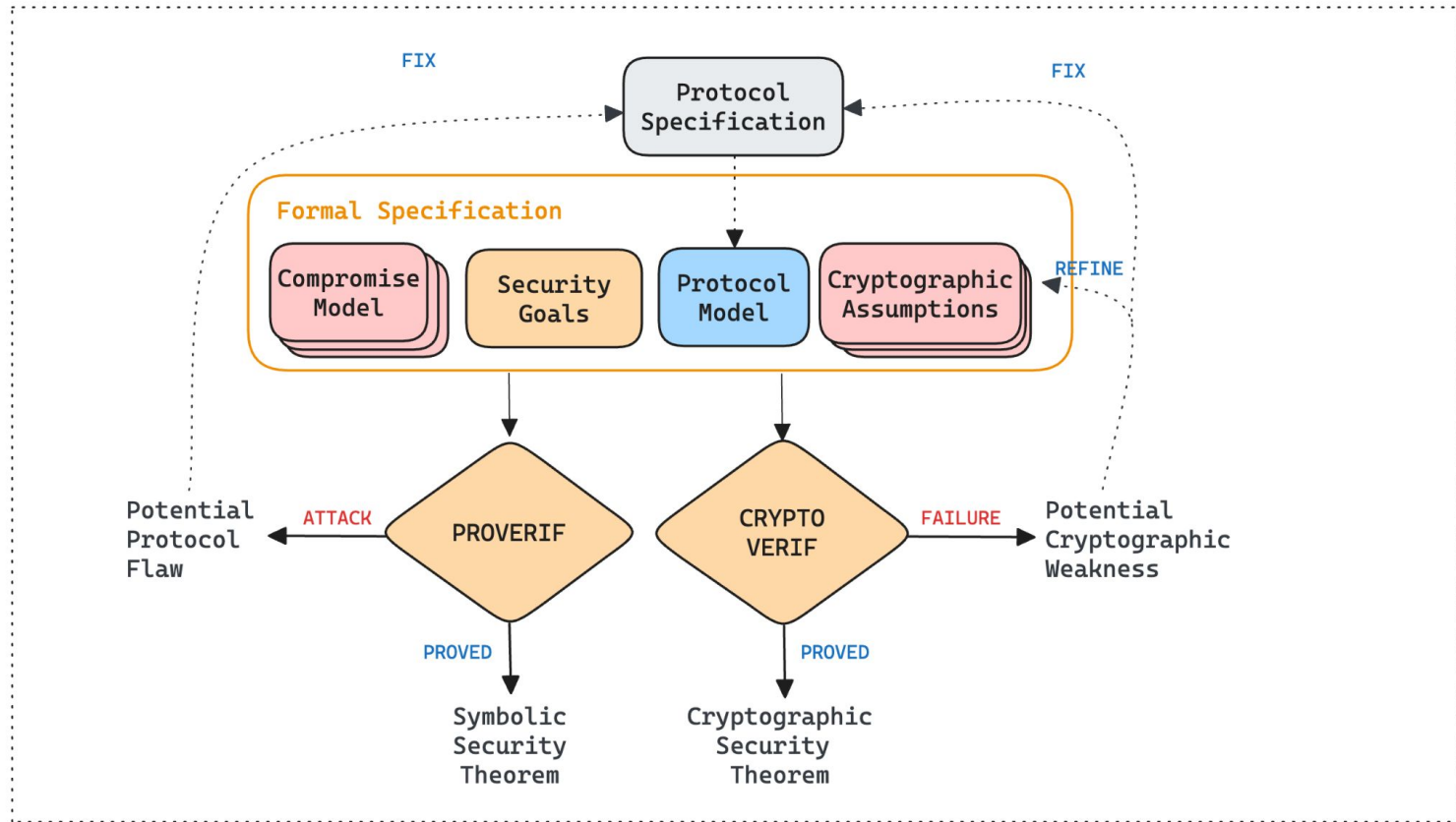
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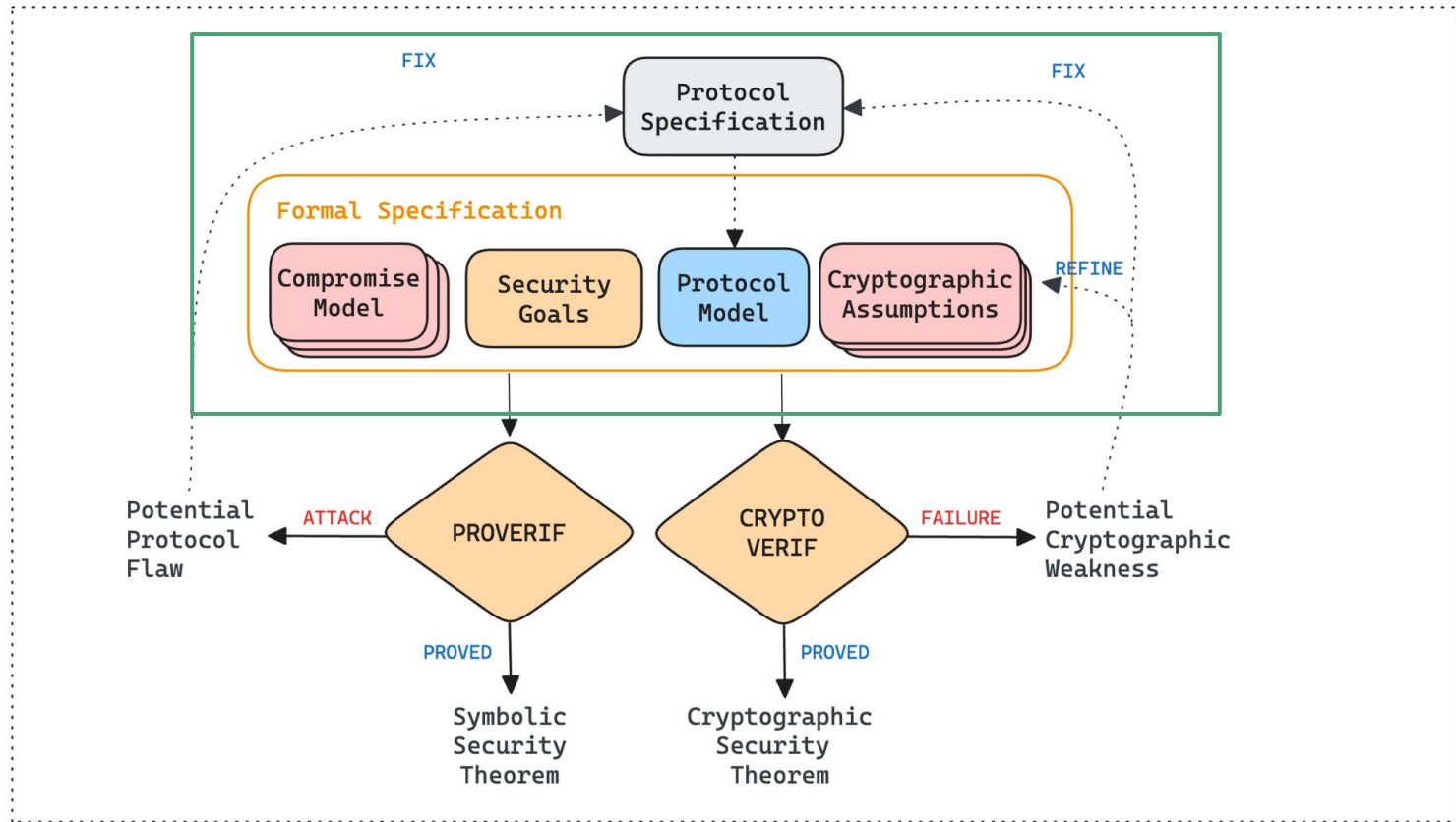
Our Formal Verification Methodology | Security



Our Formal Verification Methodology | Security



Our Formal Verification Methodology | Security



hax: ongoing projects

- **Rust Core:** an annotated version of the Rust Core library
- **Backends:** new backends for Lean, EasyCrypt, ProVerif
- **Verified**
 - **PQ Crypto:** verified Rust code for Kyber/ML-KEM, ...
 - **OS Modules:** verified kernel code for RIOT-OS
 - **Protocols:** verified code for EDHOC, MLS, TLS 1.3, ...
 - **Contracts:** verified canisters for Internet Computer

HAX

A Usable Tool for
Verification



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